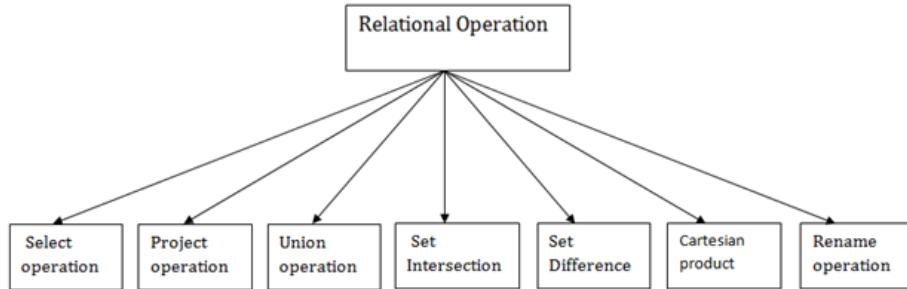


## Relational Algebra

Relational algebra is a procedural query language. It gives a step by step process to obtain the result of the query. It uses operators to perform queries.

### Types of Relational operation



#### 1. Select Operation:

- The select operation selects tuples that satisfy a given predicate.
- It is denoted by sigma ( $\sigma$ ).

Notation:  $\sigma p(r)$

#### Where:

- $\sigma$  is used for selection prediction
- $r$  is used for relation
- $p$  is used as a propositional logic formula which may use connectors like: AND OR and NOT. These relational can use as relational operators like  $=, \neq, \geq, <, >, \leq$ .

For example: LOAN Relation



BRANCH_NAME	LOAN_NO	AMOUNT
Downtown	L-17	1000
Redwood	L-23	2000
Perryride	L-15	1500
Downtown	L-14	1500
Mianus	L-13	500
Roundhill	L-11	900
Perryride	L-16	1300

**Input:**

```
σ BRANCH_NAME="perryride" (LOAN)
```

**Output:**

BRANCH_NAME	LOAN_NO	AMOUNT
Perryride	L-15	1500
Perryride	L-16	1300

## 2. Project Operation:

- o This operation shows the list of those attributes that we wish to appear in the result. Rest of the attributes are eliminated from the table.
- o It is denoted by  $\Pi$ .

```
Notation:  $\Pi A_1, A_2, A_n (r)$ 
```

**Where**

**A1, A2, A3** is used as an attribute name of relation **r**.

**Example: CUSTOMER RELATION**

NAME	STREET	CITY
Jones	Main	Harrison
Smith	North	Rye
Hays	Main	Harrison
Curry	North	Rye
Johnson	Alma	Brooklyn
Brooks	Senator	Brooklyn



**Input:**

$\Pi \text{NAME, CITY} (\text{CUSTOMER})$
---

**Output:**

NAME	CITY
Jones	Harrison
Smith	Rye
Hays	Harrison
Curry	Rye
Johnson	Brooklyn
Brooks	Brooklyn

**3. Union Operation:**

- o Suppose there are two tuples R and S. The union operation contains all the tuples that are either in R or S or both in R & S.
- o It eliminates the duplicate tuples. It is denoted by  $\cup$ .

Notation: $R \cup S$
----------------------

A union operation must hold the following condition:

- o R and S must have the attribute of the same number.
- o Duplicate tuples are eliminated automatically.

**Example:****DEPOSITOR RELATION**

CUSTOMER_NAME	ACCOUNT_NO
Johnson	A-101
Smith	A-121
Mayes	A-321
Turner	A-176
Johnson	A-273
Jones	A-472
Lindsay	A-284

**BORROW RELATION**

CUSTOMER_NAME	LOAN_NO
Jones	L-17
Smith	L-23
Hayes	L-15
Jackson	L-14
Curry	L-93
Smith	L-11
Williams	L-17

**Input:**

$$\Pi \text{ CUSTOMER\_NAME} (\text{BORROW}) \cup \Pi \text{ CUSTOMER\_NAME} (\text{DEPOSITOR})$$
**Output:**

CUSTOMER_NAME
Johnson
Smith
Hayes
Turner
Jones
Lindsay
Jackson
Curry
Williams
Mayes

#### 4. Set Intersection:

- o Suppose there are two tuples R and S. The set intersection operation contains all tuples that are in both R & S.
- o It is denoted by intersection  $\cap$ .

Notation:  $R \cap S$

**Example:** Using the above DEPOSITOR table and BORROW table

**Input:**

$$\Pi \text{ CUSTOMER\_NAME} (\text{BORROW}) \cap \Pi \text{ CUSTOMER\_NAME} (\text{DEPOSITOR})$$

**Output:**

CUSTOMER_NAME
Smith
Jones

**5. Set Difference:**

- o Suppose there are two tuples R and S. The set intersection operation contains all tuples that are in R but not in S.
- o It is denoted by intersection minus (-).

Notation:  $R - S$

**Example:** Using the above DEPOSITOR table and BORROW table

**Input:**

$\prod \text{CUSTOMER\_NAME} (\text{BORROW}) - \prod \text{CUSTOMER\_NAME} (\text{DEPOSITOR})$

**Output:**

CUSTOMER_NAME
Jackson
Hayes
Willians
Curry

**6. Cartesian product**

- o The Cartesian product is used to combine each row in one table with each row in the other table. It is also known as a cross product.
- o It is denoted by X.

Notation:  $E \times D$

**Example:**

**EMPLOYEE**

EMP_ID	EMP_NAME	EMP_DEPT
1	Smith	A
2	Harry	C



3	John	B
---	------	---

**DEPARTMENT**

DEPT_NO	DEPT_NAME
A	Marketing
B	Sales
C	Legal

**Input:**

EMPLOYEE X DEPARTMENT

**Output:**

EMP_ID	EMP_NAME	EMP_DEPT	DEPT_NO	DEPT_NAME
1	Smith	A	A	Marketing
1	Smith	A	B	Sales
1	Smith	A	C	Legal
2	Harry	C	A	Marketing
2	Harry	C	B	Sales
2	Harry	C	C	Legal
3	John	B	A	Marketing
3	John	B	B	Sales
3	John	B	C	Legal

## 7. Rename Operation:

The rename operation is used to rename the output relation. It is denoted by **ρho** ( $\rho$ ).

**Example:** We can use the rename operator to rename STUDENT relation to STUDENT1.

ρ(STUDENT1, STUDENT)

Note: Apart from these common operations Relational algebra can be used in Join operations.

 [Youtube For Videos Join Our Youtube Channel: Join Now](#)

## Help Others, Please Share



## Learn Latest Tutorials

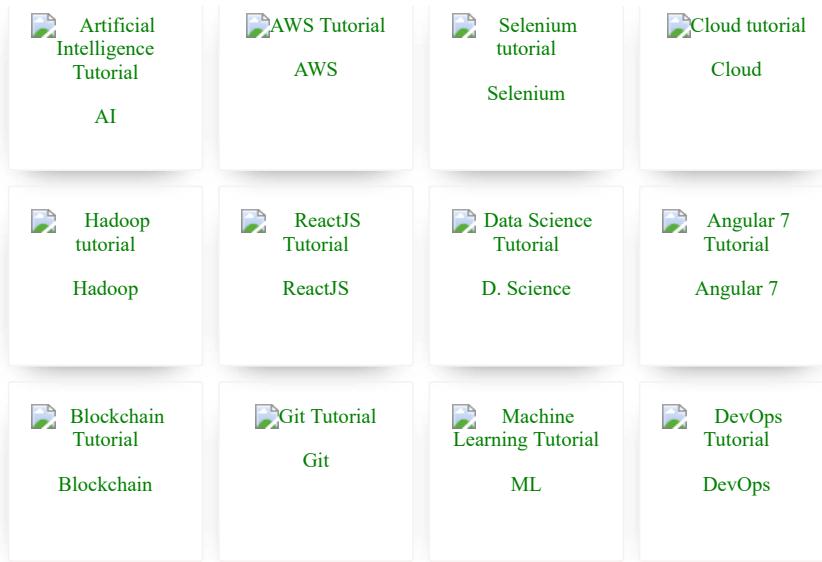
 PostgreSQL tutorial PostgreSQL	 Apache Solr Tutorial Solr	 MongoDB tutorial MongoDB	 Gimp Tutorial Gimp
 Verilog Tutorial Verilog	 Teradata Tutorial Teradata	 PhoneGap Tutorial PhoneGap	 Gmail Tutorial Gmail
 Vue.js Tutorial Vue.js	 PLC tutorial PLC	 Adobe Illustrator Tutorial Illustrator	

## Preparation

 Aptitude Aptitude	 Logical Reasoning Reasoning	 Verbal Ability Verbal A.	 Interview Questions Interview
 Company Interview Questions Company			

## Trending Technologies





## B.Tech / MCA

